

Mutability and advanced control structures

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 - A machine he devised to abstract “computation through a purely mechanical process”.

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- Church, 1936: λ -calculus.
 - Functions can be defined.
 - Functions can be called.
- Turing, 1936 (later): Turing machines.
 - A machine he devised to abstract “computation through a purely mechanical process”.
 - Consists of a memory tape, a head, a machine state, and a function that makes decisions.

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- We have been using a functional programming style, in spirit of λ -calculus.
- Computers nowadays mostly are in the style of the von Neumann architecture.
- Most programming languages have a way to mutate data.
- Most data structures and algorithms also mutate data.
- We would also like to mutate data, then.

Mutability

- `const` constant declaration.

¹See lecture notes for pitfalls.

Mutability

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- `let` *variable* declaration.

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Mutability

- `const` constant declaration.
- `let` *variable* declaration.
- Mutability leads to more complicated reasoning.¹. Yet if done right, is more “natural” to work with.

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Arrays

The Array

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Example

```
let arr = [1, 2, 3, 4];
```

Q: `arr[0]` = ?, `arr[1]` = ?, `arr[2]` = ?

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Example

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let arr = [1, 2, 3, 4];  
arr[0] = 1, arr[1] = 2, arr[2] = 3
```

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Example

```
let arr = [1, 2, 3, 4];
```

```
arr[0] = 1, arr[1] = 2, arr[2] = 3
```

```
Q: arr[4] = ?
```

Arrays

The Array

The array is initialized with a list of comma delimited items within square brackets. The items are 0-indexed. Elements can be accessed with the dereference operator `[]`.

Example

```
let arr = [1, 2, 3, 4];  
arr[0] = 1, arr[1] = 2, arr[2] = 3  
arr[4] = undefined
```

Arrays

Insertion

We can insert or modify any element by simply dereferencing it and setting it.

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Example

```
let arr = [1, 2, 3, 4];  
arr[3] = 3; arr[4] = 4; arr[6] = 6;
```

Q: arr = ?

Arrays

Insertion

We can insert or modify any element by simply dereferencing it and setting it.

Example

```
let arr = [1, 2, 3, 4];  
arr[3] = 3; arr[4] = 4; arr[6] = 6;  
[1, 2, 3, 3, 4, undefined, 6]
```


While loops

Definition

A while loop is made as such:

```
while (expression) {  
    statements  
}
```

This is something like `expression ? statements : undefined`, over and over again.

For loops

Definition

A for loop is made as such:

```
for (expr a; expr b; expr c) {  
    statements  
}
```

This is like

```
expr a;  
while(expr b;) {  
    expr c;  
    statements;  
}
```

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for (expr a; expr b; expr c) {  
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```

This is like

```
expr a;  
while(expr b;) {  
    expr c;  
    statements;  
}
```

Traversing arrays

```
for (let i = 0; i < array_length(arr); i = i + 1) {...}
```

Loop controls

Control statements

If you wish to escape the closest loop prematurely, use `break`. If you wish to skip one iteration immediately, use `continue`.

Question

Two nested loops:

```
while(...) {  
    while(...) {  
        //I am here  
    }  
}
```

How to break out of both loops?

Crisis

Substitution model does not work any more. We need a new way to keep track of our names!

Remember that time...

- Every new context creates a new *scope*.
- The most common context are blocks:
- Names in an inner scope inherit those defined outside of it.
- Names can be overridden by definitions in the current scope.
- You cannot go “into” an inner scope from an outer scope to retrieve definitions!
- **In conclusion:** To find what a name refers to, look at the current scope, and then outwards. Take the first one you come across.

Environment Model

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- New scopes create child environments.
 - All whole program resides in the program environment.
 - The parent is the closest *enclosing environment*.
- Duplicate names in the same environment cannot exist.
 - However, they can be overwritten depending on the statement.
- Environments are never destroyed.
- When evaluating a name, search outwards starting from the current environment. First match is returned. Otherwise, invalid name.

S9 Q1

```
function change(x, new_value) {  
    x = new_value;  
}  
let x = 0;  
change(x, 1);
```

What is the value of `x` after evaluation?

S9 Q1

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`x = 0.`

S9 Q2

Write `d_filter(xs)` that acts like `filter` but modifies the list `xs` in place.

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```
function d_filter(pred, xs) {
```

```
}
```

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```
function d_filter(pred, xs) {  
  if (is_null(xs)) {  
    return xs;  
  }  
  
}
```

S9 Q2

Write `d_filter(xs)` that acts like `filter` but modifies the list `xs` in place.

```
function d_filter(pred, xs) {  
  if (is_null(xs)) {  
    return xs;  
  }  
  else if (pred(head(xs))) {  
    set_tail(xs, d_filter(pred, tail(xs)));  
    return xs;  
  }  
  
}
```

S9 Q2

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function d_filter(pred, xs) {
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  }
  else if (pred(head(xs))) {
    set_tail(xs, d_filter(pred, tail(xs)));
    return xs;
  }
  else {
    return d_filter(pred, tail(xs));
  }
}
```

S9 Q3

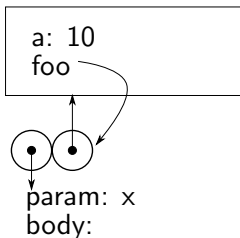
Draw the environment at the breakpoints.

```
let a = 10;
function foo(x) {
  let b = 0;
  function goo(x) {
    let a = 30;
  }
  if (x <= 2) {
    a = a + x;
    b = b + x;
    // Breakpoint #4
  } else {
    // Breakpoint #3
    goo(x - 1);
  }
  a = a + x;
  b = b + x;
  // Breakpoint #2
  goo(3);
}
// Breakpoint #1
foo(1);
// Breakpoint #5
```


S9 Q3

Breakpoint 1

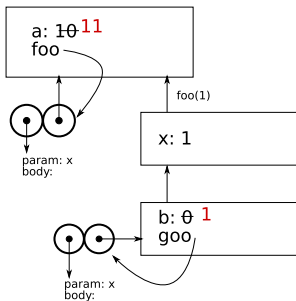
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let a = 10;
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    let a = 30;
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  if (x <= 2) {
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    b = b + x;
  } else {
    goo(x - 1);
  }
  a = a + x;
  b = b + x;
  goo(3);
}
// Breakpoint #1
foo(1);
```



S9 Q3

Breakpoint 2

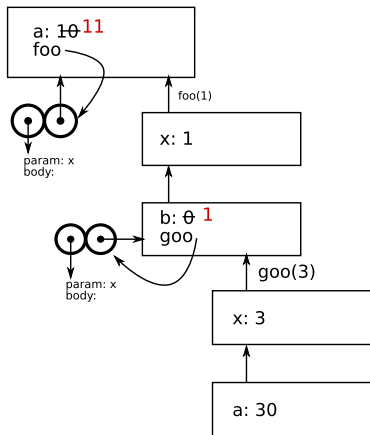
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  function goo(x) {
    let a = 30;
  }
  if (x <= 2) {
    a = a + x;
    b = b + x;
  } else {
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  }
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  b = b + x;
  // Breakpoint #2
  goo(3);
}
foo(1);
```



S9 Q3

Breakpoint 3

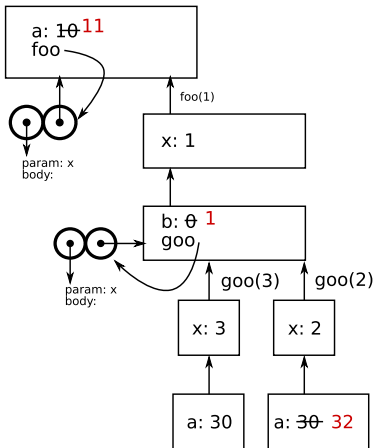
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  }
  if (x <= 2) {
    a = a + x;
    b = b + x;
  } else {
    // Breakpoint #3
    goo(x - 1);
  }
  a = a + x;
  b = b + x;
  goo(3);
}
foo(1);
```



S9 Q3

Breakpoint 4

```
let a = 10;
function foo(x) {
  let b = 0;
  function goo(x) {
    let a = 30;
  }
  if (x <= 2) {
    a = a + x;
    b = b + x;
    // Breakpoint #4
  } else {
    goo(x - 1);
  }
  a = a + x;
  b = b + x;
  goo(3);
}
foo(1);
```



Q1

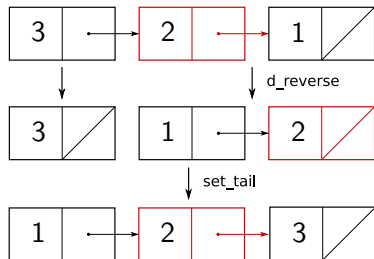
Write `d_reverse(xs)`, just like `d_filter` you did just now.

```
const L = list(1, 2, 3, 4, 5, 6);  
d_reverse(L); // returns [6, [5, [4, [3, [2, [1, null]]]]]]
```

Q1

Write `d_reverse(xs)`, just like `d_filter` you did just now.

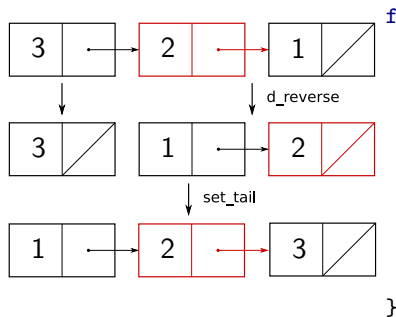
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const L = list(1, 2, 3, 4, 5, 6);  
d_reverse(L); // returns [6, [5, [4, [3, [2, [1, null]]]]]]
```



```
function d_reverse(xs) {  
  if (is_null(xs) || is_null(tail(xs)))  
    return xs;  
  }  
  else {  
    const temp = d_reverse(tail(xs));  
    set_tail(tail(xs), xs);  
    set_tail(xs, null);  
    return temp;  
  }  
}
```

Q2

Ex. 3.16

Consider the following function that counts the number of pairs in a list. Give examples of lists made of 3 pairs that will cause the function to return 3, 4, 7, or never return at all.

```
function count_pairs(x) {  
  if (!is_pair(x)) {  
    return 0;  
  } else {  
    return 1 + count_pairs(head(x)) + count_pairs(tail(x));  
  }  
}
```


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}
```

```
// returns 3  
const three = list(1, 2, 3);  
count_pairs(three);
```

Q2

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  } else {  
    return 1 + count_pairs(head(x)) + count_pairs(tail(x));  
  }  
}
```

```
// returns 3  
const three = list(1, 2, 3);  
count_pairs(three);
```

```
// infinite loop  
const loop = list(1, 2, 3);  
set_tail(tail(tail(loop)), loop);  
count_pairs(loop);
```

```
// returns 4  
const four_a = pair(null, null);  
const four_b = pair(four_a, four_a);  
const four = pair(four_b, null);  
count_pairs(four);
```

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  } else {  
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  }  
}
```

```
// returns 3  
const three = list(1, 2, 3);  
count_pairs(three);
```

```
// infinite loop  
const loop = list(1, 2, 3);  
set_tail(tail(tail(loop)), loop);  
count_pairs(loop);
```

```
// returns 4  
const four_a = pair(null, null);  
const four_b = pair(four_a, four_a);  
const four = pair(four_b, null);  
count_pairs(four);
```

```
// returns 7  
const seven_a = pair(null, null);  
const seven_b = pair(seven_a, seven_a);  
const seven = pair(seven_b, seven_b);  
count_pairs(seven);
```